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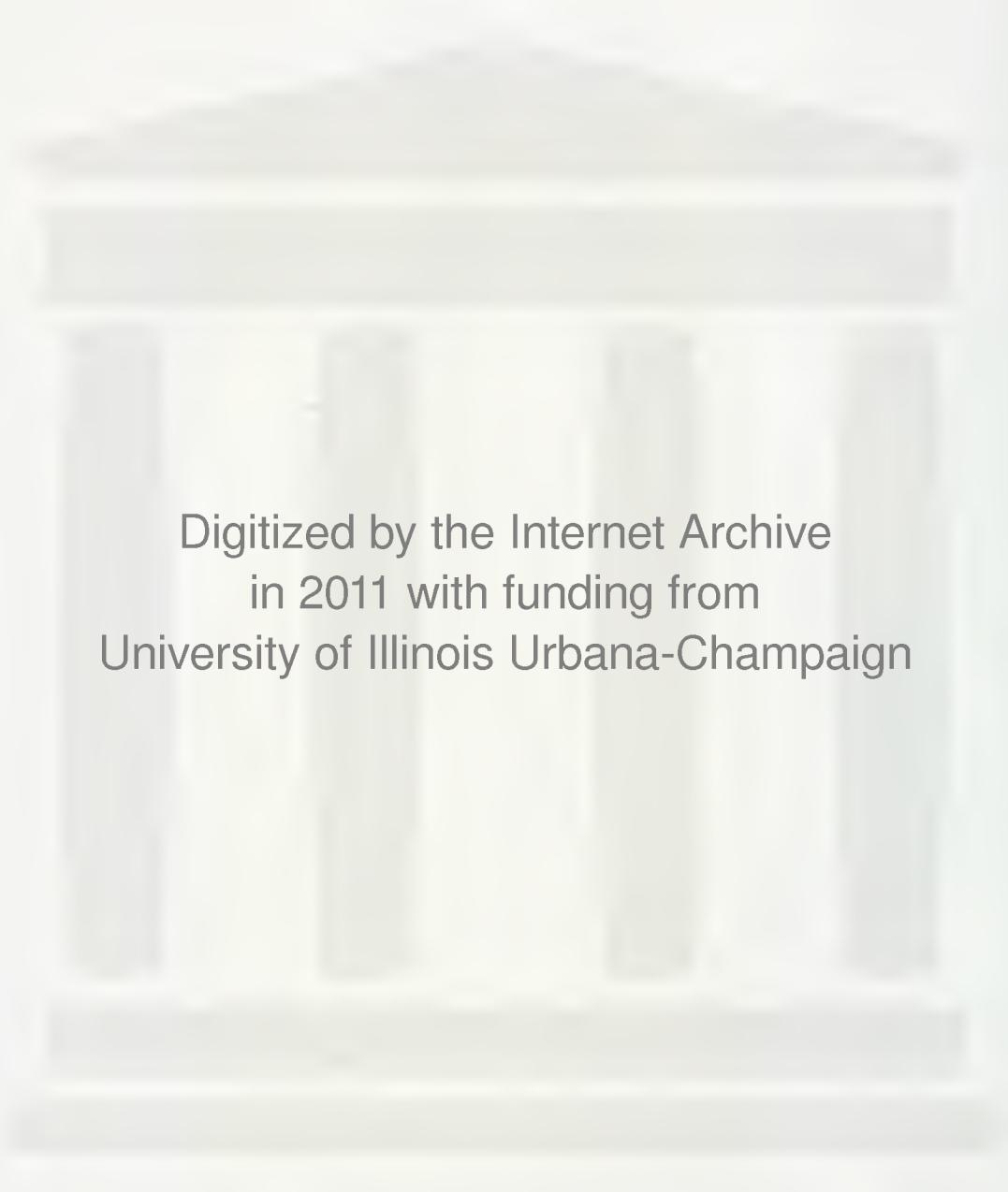
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## Faculty Working Papers

CONVEX OPTIMIZATION AND LAGRANGE MULTIPLIERS

Charles E. Blair

#407

**College of Commerce and Business Administration**  
**University of Illinois at Urbana-Champaign**



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APPENDIX

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Convex Optimization and Lagrange Multipliers

by

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June 8, 1977

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Abstract

We show how the duality theorem of linear programming can be used to  
prove several results on general convex optimization.



Let  $f, g_1, \dots, g_k$  be convex functions defined on a convex subset  $S$  of a vector space. Let  $T = \{x \in S \mid g_i(x) \leq 0 \ 1 \leq i \leq k\}$ . We assume throughout that  $T$  is non-empty.

We use linear programming theory to explore the relationship between the problem of minimizing  $f(x)$   $x \in T$  and the Lagrange dual problem of minimizing  $f(x) + \sum \lambda_i g_i(x)$   $x \in S$  for suitable  $\lambda_i \geq 0$ . This is motivated by the work of Duffin [1, 2, 3]. The main tool we shall need is a version of the duality theorem of linear programming [5, theorems 1.1.9 and 1.7.13].

Lemma: Let  $A = \{x \mid Bx \geq b \text{ and } Cx = c\}$ . If  $A$  is empty, there are  $U \geq 0$  and  $V$  such that  $U^t B + V^t C = \vec{0}$  and  $Ub + Vc > 0$ . If  $A$  is non-empty and every  $x \in A$  satisfies  $dx \geq e$ , then there are  $U \geq 0$  and  $V$  such that  $U^t B + V^t C = d$  and  $Ub + Vc \geq e$ .

We will not require separating hyperplane theorems or results from semi-infinite programming.

Theorem 1:  $f(x) \geq L$  for every  $x \in T$  if and only if for every finite FCS there are  $\lambda_1, \dots, \lambda_k$  such that  $f(x) + \sum \lambda_i g_i(x) \geq L$  for every  $x \in F$ .

Proof: The "if" part is immediate. For the "only if" part it suffices to prove the result for those finite  $F$  which contain members of  $T$ . Let  $F = \{y_1, \dots, y_N\}$   $y_i \in T$ . For such  $F$  the system of equations and inequalities in unknowns  $\theta_1, \dots, \theta_N$

$$(D) \quad \begin{aligned} & \sum_{i=1}^N \theta_i = 1 \\ & \sum_{i=1}^N \theta_i g_j(y_i) \leq 0 \quad 1 \leq j \leq k \\ & \theta_i \geq 0 \end{aligned}$$

has the solution  $\theta_i = 1$ . By convexity, if  $\theta_i$  is a solution to (D),  $\sum \theta_i y_i \in T$ .



Since we are assuming  $f(x) \geq L$  for  $x \in T$ , every solution to (D) must satisfy  $\sum_{i=1}^k \lambda_i f(y_i) \geq L$ . By the lemma there are  $\lambda_i \geq 0$  and  $\gamma$  such that

$$\gamma + \sum_{j=1}^k \lambda_j (-g_j(y_i)) \leq f(y_i) \quad 1 \leq i \leq N$$

$$\gamma \geq L$$

so  $f(y_i) + \sum_{j=1}^k \lambda_j g_j(y_i) \geq L$  for every  $y_i \in F$ . Q.E.D.

Theorem 1 may be used to prove many of the standard results on convex optimization. As an example we prove the Kuhn-Tucker theorem.

Corollary: Suppose  $f(x) \geq L$  for  $x \in T$  and that there is a  $y$  for which  $g_i(y) < 0$   $1 \leq i \leq k$ . Then there are  $\lambda_1, \dots, \lambda_k$  such that  $f(x) + \sum_{i=1}^k \lambda_i g_i(x) \geq L$  for  $x \in S$ .

Proof: Let  $\delta = \max \{g_i(y)\}$ . For  $x \in S$  let  $A_x = \{(\lambda_1, \dots, \lambda_k) \mid \lambda_i \geq 0, f(x) + \sum_{i=1}^k \lambda_i g_i(x) \geq L, \text{ and } -\delta(\sum \lambda_i) \leq f(y) - L\}$ . For each  $x \in S$ ,  $A_x$  is compact. If  $H \cap S$  is finite we may use theorem 1 with  $F = H \cup \{y\}$  to show  $\bigcap_{x \in H} A_x$  is non-empty. Therefore,  $\bigcap_{x \in S} A_x$  is non-empty, so suitable Lagrange multipliers exist. Q.E.D.

Arguments of this kind can also be used to give information about when "duality gaps" occur.

Theorem 2. There are no  $\lambda_i$  such that  $f(x) + \sum_{i=1}^k \lambda_i g_i(x) \geq L$  for every  $x \in S$  if and only if, for every  $N > 0$ , there is an  $x \in S$  such that  $f(x) < L - N(\max_{1 \leq i \leq k} g_i(x))$ .

Proof: Take any  $\lambda_i \geq 0$  and suppose  $x$  exists with the desired property for  $N = \sum \lambda_i$ . Then  $f(x) + \sum_{i=1}^k \lambda_i g_i(x) \leq f(x) + (\max_{i=1}^k g_i(x)) (\sum \lambda_i) < L$ . So no suitable  $\lambda_i$  exist.

Conversely, suppose there are no suitable  $\lambda_i$ . For any  $N > 0$  and any  $F = \{y_1, \dots, y_M\} \subset S$  consider the linear system in unknowns  $\lambda_1, \dots, \lambda_k$ :



$$f(y_i) + \sum_{j=1}^k \lambda_j g_j(y_i) \geq L \quad y_i \in F$$

$$(F;N) \quad \sum \lambda_i \leq N$$

$$\lambda_i \geq 0$$

If, for some  $N$ ,  $(F;N)$  had a solution for every finite FCS, a compactness argument similar to that in the corollary to Theorem 1 would yield suitable multipliers  $\lambda_i$ . Since we are assuming such  $\lambda_i$  do not exist, it must be that for every  $N > 0$  there is an  $F$  such that  $(F;N)$  has no solution. By the lemma, if  $(F;N)$  has no solution, there are  $\theta_1, \dots, \theta_M \geq 0$  and  $\gamma > 0$  such that

$$\sum_{i=1}^M \theta_i g_j(y_i) - \gamma \leq 0, \quad 1 \leq j \leq k$$

and  $\sum \theta_i (L - f(y_i)) + \gamma(-N) > 0$ . By scaling, we may assume  $\sum_{i=1}^M \theta_i = 1$ , so that

$$\sum_{i=1}^M \theta_i g_j(y_i) \leq \gamma < \frac{1}{N} (L - \sum \theta_i f(y_i)).$$

If we take  $x = \sum \theta_i y_i$ ,  $\gamma \geq g_j(x)$ ,  $1 \leq j \leq k$  and  $f(x) < L - N\gamma$  follow by convexity of  $f$  and  $g_j$ . Q.E.D.

Corollary: (Compare [1], cor. 5; [2], thm. 3): Let  $h(\epsilon) = \inf\{f(x) \mid g_j(x) \leq \epsilon, 1 \leq j \leq k\}$ . If there are  $\delta > 0$ ,  $L$  such that  $h(x) > L$  for  $0 < x \leq \delta$ , then there are  $\lambda_1, \dots, \lambda_k$  such that  $f(x) + \sum \lambda_i g_i(x) \geq L \quad x \in S$ .

Proof: If there is an  $x$  for which  $g_j(x) < 0$  the existence of suitable  $\lambda_i$  follows from the Kuhn-Tucker theorem,\* so we assume this is not the case.  $h$  is a convex monotone function which, on our assumptions, is defined only for non-negative arguments. For  $\epsilon \geq \delta$ ,  $h(\epsilon) - L > h(\epsilon) - h(0) \geq \epsilon \left(\frac{1}{\delta}\right) (h(\delta) - h(0))$ . Hence for  $x \in S$ ,  $f(x) - L \geq (\max g_i(x)) (\min 0, \frac{1}{\delta} (h(\delta) - h(0)))$ . (Note that our assumptions

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\* We use the Kuhn-Tucker theorem for brevity. The result could be proved from Theorem 2 alone.





imply  $\max g_j(x) \geq 0$ .) Since the condition given by Theorem 2 fails for  $N = \max(0, \frac{1}{\delta}(h(0) - h(\delta)))$ , suitable  $\lambda_i$  exist.) Q.E.D.

Finally, we use a variation of these techniques to strengthen a recent result of Duffin and Jeroslow [4].

Theorem 3: Let  $S = \mathbb{R}^n$ . Assume that for  $\lambda_i > 0$ ,  $f(x) + \sum \lambda_i g_i(x) \geq L$ ,  $x \in S$ . Then there are affine functions  $h_i(x) = a_i x + b_i$  ( $a_i \in \mathbb{R}^n$ ,  $b_i \in \mathbb{R}$ ) such that  $h_i(x) \leq g_i(x)$  and  $f(x) + \sum \lambda_i h_i(x) \geq L$ ,  $x \in S$ .

Proof: For  $y \in S$ , let  $T_y = \{(h_1, \dots, h_k) \mid h_i \text{ affine, } h_i(y) \leq g_i(y), \text{ and } f(y) + \sum \lambda_i h_i(y) \geq L\}$ .

We identify each function  $h_i(x) = a_i x + b_i$  with the ordered pair  $(a_i, b_i) \in \mathbb{R}^{n+1}$ . Thus,  $T_y$  is a closed subset of  $\mathbb{R}^{(n+1)k}$ . We first show that, for any finite FCS  $\bigcap_{y \in F} T_y$  is non-empty. A member of  $\bigcap_{y \in F} T_y$  would be a solution to the linear inequality system in unknowns  $a_1, \dots, a_k; b_1, \dots, b_k$ .

$$\begin{aligned} a_i y + b_i &\leq g_i(y) & 1 \leq i \leq k; & y \in F \\ \sum_{i=1}^k \lambda_i (a_i y + b_i) &\geq M - f(y) & y \in F \end{aligned} \quad (E)$$

By the lemma, (E) has no solution only if there are scalars  $W_{iy}, V_y \geq 0$  such that

$$\begin{aligned} (i) \quad \sum_{y \in F} W_{iy} y &= \lambda_i \sum_{y \in F} V_y y, & 1 \leq i \leq k \\ (ii) \quad \sum_{y \in F} W_{iy} &= \lambda_i \sum_{y \in F} V_y, & 1 \leq i \leq k \\ (iii) \quad \sum_{\substack{1 \leq i \leq k \\ y \in F}} W_{iy} g_i(y) &< \sum_{y \in F} V_y (M - f(y)). \end{aligned}$$

If there were  $W, V$  satisfying (i)-(iii) we could set  $V'_y = V_y / \sum_{iy} W_{iy}$  and

$W'_{iy} = W_{iy} / \lambda_i \sum_{iy} W_{iy}$  so that (i)-(iii) would be satisfied and, by (ii),  $\sum_{y \in F} V'_y = \sum_{y \in F} W'_{iy} = 1$ ,  $1 \leq i \leq k$ . Condition (i) becomes  $\sum W'_{iy} y = \sum V'_y y$ . Condition (iii) becomes  $\sum_{y \in F} V'_y f(y) + \sum_{y \in F} \lambda_i ( \sum_{iy} W'_{iy} g_i(y) ) < M$ . If  $z = \sum W'_{iy} y = \sum V'_y y$  this implies, by convexity,  $f(z) + \sum \lambda_i g_i(z) < M$



which contradicts our assumption about the  $\lambda_i$ . Therefore (E) has solutions for every finite F.

To complete the proof we must show  $\bigcap_{y \in S} T_y$  is non-empty. Let  $e_j$  = jth unit vector. We show that if F contains  $\pm e_j$   $1 \leq j \leq n$  and the zero vector, then  $\bigcap_{y \in F} T_y$  is bounded. Since each  $T_y$  is closed, compactness yields the desired result. For  $1 \leq i \leq k$  we must have  $h_i(\vec{0}) = b_{i-} \leq g_i(0)$ ,  $f(\vec{0}) + \sum \lambda_i b_{i-} > L$ ,  $h_i(e_j) = (j\text{th component of } a_i) + b_{i-} \leq g_i(e_j)$ , and  $h_i(-e_j) \leq g_i(-e_j)$ . Since all the  $\lambda_i$  are positive this implies bounds on  $a_i$ ,  $b_{i-}$ . Q.E.D.

Our proof of Theorem 3 works for any convex  $S \subset \mathbb{R}^n$  which includes  $\pm e_j$  and  $\vec{0}$ . By suitable translations, this implies the results for any fully dimensional convex  $S \subset \mathbb{R}^n$ . Further modifications yield the result for arbitrary convex  $S \subset \mathbb{R}^n$ .

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